

Math 821, Spring 2013, Lecture 13

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1 SubHopf Algebra Results

Example. Last time we showed that $T(x) = \mathbb{1} - B_+(\frac{1}{T(x)})$.

$$T(x) = \mathbb{1} - \bullet \cdot x - \downarrow x^2 - \left(\begin{array}{c} \bullet \\ | \\ \bullet \\ | \\ \bullet \end{array} + \begin{array}{c} \bullet \\ \diagup \\ \bullet \end{array} \right) x^3 - \left(\begin{array}{c} \bullet \\ | \\ \bullet \\ | \\ \bullet \end{array} + \begin{array}{c} \bullet \\ \diagup \\ \bullet \end{array} + 2 \begin{array}{c} \bullet \\ \diagup \\ \bullet \\ \diagdown \\ \bullet \end{array} + \begin{array}{c} \bullet \\ \diagup \\ \bullet \\ \diagdown \\ \bullet \end{array} \right) x^4 + O(x^5).$$

Define $t_0 = \mathbb{1}, n > 0, t_n = -[x^n]T(x)$. Now what is $\Delta(t_i)$?

$$\Delta(t_0) = \mathbb{1} \otimes \mathbb{1} = t_0 \otimes t_0$$

$$\Delta(t_1) = \mathbb{1} \otimes \bullet + \bullet \otimes \mathbb{1} = t_0 \otimes t_1 + t_1 \otimes t_0$$

$$\Delta(t_2) = \mathbb{1} \otimes t_2 + t_2 \otimes \mathbb{1} + \bullet \otimes \bullet = t_0 \otimes t_2 + t_2 \otimes t_0 + t_1 \otimes t_1$$

$$\begin{aligned} \Delta(t_3) &= t_0 \otimes t_3 + t_3 \otimes t_0 + 3 \bullet \otimes \downarrow + \downarrow \otimes \bullet + \bullet \otimes \bullet \\ &= t_0 \otimes t_3 + t_3 \otimes t_0 + 3t_1 \otimes t_2 + (t_2 + t_1^2) \otimes t_1 \end{aligned}$$

$$\Delta(t_4) = \Delta \left(\begin{array}{c} \bullet \\ | \\ \bullet \\ | \\ \bullet \end{array} + \begin{array}{c} \bullet \\ \diagup \\ \bullet \end{array} + 2 \begin{array}{c} \bullet \\ \diagup \\ \bullet \\ \diagdown \\ \bullet \end{array} + \begin{array}{c} \bullet \\ \diagup \\ \bullet \\ \diagdown \\ \bullet \end{array} \right)$$

$$\begin{aligned} &= t_4 \otimes t_0 + t_0 \otimes t_4 + \bullet \otimes 5 \downarrow + \bullet \otimes 5 \begin{array}{c} \bullet \\ \diagup \\ \bullet \end{array} + 3 \downarrow \otimes \downarrow + 6 \bullet \otimes \downarrow + \downarrow \otimes \bullet \\ &\quad + \begin{array}{c} \bullet \\ \diagup \\ \bullet \end{array} \otimes \bullet + \downarrow \otimes \bullet \otimes 2 \bullet + \bullet \otimes \bullet \otimes \bullet \\ &= t_4 \otimes t_0 + t_0 \otimes t_4 + 5t_1 \otimes t_3 + 3(t_2 + 2t_1^2) \otimes t_2 + (t_3 + 2t_2t_1 + t_1^3) \otimes t_1 \end{aligned}$$

What we are observing is if A be the algebra generated by the t_i then $\Delta(t_i) \subseteq A \otimes A$. So A is not just a subalgebra of the Connes-Kreimer Hopf algebra; it is also a subHopf algebra.

Theorem 1 [1] *Let H be a graded connected Hopf algebra which is either free*

(words over generators) or free-commutative (polynomial algebra over generators) as an algebra. Let $(B_+^n)_{n=1}^\infty$ be a family of Hochschild 1-cocycles. Then

$$T(x) = \mathbb{1} + \sum_{n=1}^{\infty} x^n w_n B_+^n(T(x)^{n+1}),$$

where the w_n are from k . This has a unique solution given recursively and $\Delta t_n = \sum_{j=0}^n P_{n,j}(t_0, \dots, t_{n-j}) \otimes t_j$, where $T(x) = \sum_{n=0}^{\infty} t_n x^n$ and $P_{n,j}$ is a polynomial.

But this isn't quite satisfactory because the specification is rather special. If we only have a B_+ of weight 1, here's a nice theorem:

Theorem 2 [2] Let $P = \sum_{n=0}^{\infty} p_n x^n$ be a formal power series with $p_0 = 1$, then $T(x) = x B_+(P(T(x)))$ has a unique solution given recursively and the following are equivalent:

1. The algebra generated by the t_i ($T(x) = \sum_{n=0}^{\infty} t_n x^n$) is a subHopf algebra.
2. $\exists(\alpha, \beta) \in \mathbb{Q}^2$ such that $(1 - \alpha\beta x)P'(x) = \alpha P(x)$.
3. $\exists(\alpha, \beta) \in \mathbb{Q}^2$ such that
 - (a) $P(x) = 1$ if $\alpha = 0$.
 - (b) $P(x) = e^{\alpha x}$ if $\beta = 0, \alpha \neq 0$.
 - (c) $P(x) = (1 - \alpha\beta x)^{-\frac{1}{\beta}}$ else.

Together

Theorem 3 [3] Suppose

$$T(x) = \sum_{j \in J} x^j \beta^j + (P_j(T(x))) \quad (*)$$

with $J \subseteq \{1, 2, \dots\}$, $P_j(0) = 1$, P_j formal power series, and suppose the coefficients of the solution $T(x)$ form a subHopf algebra. Then one of the following holds:

1. $\exists \lambda, \mu \in \mathbb{Q}$ such that $(*)$ is

$$T(x) = \sum_{j \in J} x^j B_+^j((1 - \mu T(x) Q(T(x))^j),$$

where

$$Q(h) = \begin{cases} (1 - \mu\lambda)^{\frac{\lambda}{\mu}} & \text{if } \mu \neq 0 \\ e^{\lambda h} & \text{if } \mu = 0. \end{cases}$$

2. $\exists m \geq 0$ and $\alpha \in \mathbb{Q}, \alpha \neq 0$ such that (*) is

$$T(x) = \sum_{\substack{j \in J \\ m \vdots j}} x B_+^j (1 + \alpha T(x)) + \sum_{\substack{j \in J \\ m \nmid j}} x B_+^j (\mathbb{1})$$

Similar results hold for specification which are systems.

Let's prove part of Foissy's 2007 result. We'll do (1) \Rightarrow (2) \Rightarrow (3). The proof of (3) \Rightarrow (1) goes through the plane version and involves looking at reductions on the pairs (α, β) .

Proof.

(2) \Rightarrow (3) Solve the differential equation. If $\alpha = 0$ the differential equation becomes $P'(x) = 0$, so P is a constant and so by the normalization $P(x) = 1$.

Assume $\alpha \neq 0$. If $\beta = 0$, $P'(x) = \alpha P(x)$ so $P(x) = e^{\alpha x}$ (and normalize this way as $P(0) = 1$).

Assume α, β are both nonzero. Then the differential equation is $(1 - \alpha\beta x) \frac{dP}{dx} = \alpha P(x)$, so

$$\frac{dP}{P(x)} = \frac{\alpha dx}{1 - \alpha\beta x},$$

so

$$\log P + C = \int \frac{dP}{P} = \int \frac{\alpha dx}{1 - \alpha\beta x} = \log(1 - \alpha\beta x).$$

But $P(0) = 1$, so $0 + C = 0$, so $C = 0$. Therefore $\log P = \log(1 - \alpha\beta x)^{-\frac{1}{\beta}}$, so $P(x) = (1 - \alpha\beta x)^{-\frac{1}{\beta}}$.

(1) \Rightarrow (2) Let A be the algebra generated by the coefficients of T . First note that if $P(x) = 1$ then $T(x) = x$ and all is true. So from now on assume $P(x)$ has a constant term.

Suppose $p_n \neq 0, n \geq 2, n$ minimal, then $t_1 = t_2 = \dots = t_n = 0$ but

$t_{n+1} = p_n B_+(\bullet^n)$, but $B_+(\bullet^n) = \underbrace{\bullet \bullet \dots \bullet}_{n \text{ times}}$. But by hypothesis we have a

subHopf algebra, so $\Delta \left(\underbrace{\bullet \bullet \dots \bullet}_{n \text{ times}} \right) \subseteq A \otimes A$. But A contains no trees of

size $2 - n$, which is a contradiction since

$$\Delta \left(\underbrace{\bullet \bullet \dots \bullet}_{n \text{ times}} \right) = \underbrace{\bullet \bullet \dots \bullet}_{n \text{ times}} \otimes \mathbb{1} + \sum_{k=0}^n \binom{n}{k} \bullet^k \otimes \underbrace{\bullet \bullet \dots \bullet}_{n-k \text{ times}}.$$

So we need to have $t_{n-k} \neq 0, \forall 0 \leq k \leq n$, so $p_1 \neq 0$. As a further consequence there is a tree of every size in A , because $0 \neq p_1 B_+(t_n) \in$

H_{n+1} appears in t_{n+1} , where H is the Connes-Kreimer Hopf algebra $/\mathbb{Q}$. Let $Z : H \rightarrow \mathbb{Q}$ be the characteristic map of \bullet , i.e. $Z(F) = \delta_{\bullet, F}$ on forests and extended linearly. Consider $(Z \otimes id)\Delta(T(x))$. By assumption A is a subHopf algebra so $\Delta(t_n) \subseteq A \otimes A$, so $(Z \otimes id)\Delta(T(x)) \in A[[x]]$. Also observe the following

$$(Z \otimes id)(ab) = (Z \otimes id)(a)(\varepsilon \otimes id)(b) + (\varepsilon \otimes id)(a)(Z \otimes id)(b), \quad (1)$$

for $a, b \in H \otimes H$. Let's check this. It suffices to check for a, b pure tensors, i.e. $a = a_1 \otimes a_2, b = b_1 \otimes b_2$.

$$\begin{aligned} \text{LHS of (1)} &= (Z \otimes id)(a_1 b_1 \otimes a_2 b_2) \\ &= \begin{cases} a_2 b_2 & \text{if } a_1 = \bullet, b_1 = \mathbb{1} \text{ or } a_1 = \mathbb{1}, b_1 = \bullet \\ 0 & \text{otherwise.} \end{cases} \end{aligned}$$

On the RHS of (1), again if $a_1 \neq \bullet, b_1 \neq \bullet$, we get 0. If $a_1 = \bullet$ then we have $a_2 \varepsilon(b_1) b_2 + 0$, since $\varepsilon(a_1) = 0$, so for this to be nonzero we need $b_1 = \mathbb{1}$. Similarly for the second term on the RHS. (Note throughout I have pushed scalars on to the a_2, b_2 part.) This proves (1).

$$\begin{aligned} (Z \otimes id)\Delta(T(x)) &= (Z \otimes id)\Delta(xB_+(P(T(x)))) \\ &= (Z \otimes id)\Delta\left(\sum_{n=0}^{\infty} xp_n B_+(T(x)^n)\right) \\ &= \sum_{n=0}^{\infty} (Z \otimes id)\Delta B_+(T(x)^n) \\ &= \sum_{n=0}^{\infty} xp_n Z(B_+(T(x)^n)) + \sum_{n=1}^{\infty} xp_n (Z \otimes B_+)\Delta(T(x)^n). \end{aligned}$$

(Since $\Delta B_+ = B_+ \otimes \mathbb{1} + (id \otimes B_+)\Delta$, we have:)

$$= Z(T(x)) + B_+\left(\sum_{n=1}^{\infty} xp_n (Z \otimes id)\Delta(T(x)^n)\right),$$

since the constant does not affect B_+ .

$$\begin{aligned} &= t_1 + B_+\left(\sum_{n=0}^{\infty} xnp_n (\varepsilon \otimes id)\Delta(T(x))^{n-1} (Z \otimes id)\Delta(T(x))\right) \text{ by (1)} \\ &= t_1 + B_+\left(\sum_{n=0}^{\infty} xnp_n T(x)^{n-1} (Z \otimes id)\Delta(T(x))\right) \\ &= t_1 + xB_+(P'(T(x)))(Z \otimes id)\Delta(T(x)). \end{aligned}$$

Next let

$$\begin{aligned} L : H[[x]] &\longrightarrow H[[x]] \\ a &\longmapsto xB_+(P'(T(x))a). \end{aligned}$$

L increases degree, so $id - L$ is formally invertible. So the calculation above says

$$(id - L)((Z \otimes id)\Delta(T(x))) = Z(T(x))(t_1),$$

or, equivalently

$$(Z \otimes id)\Delta(T(x)) = (id - L)^{-1}(t_1) = t_1(id - L)^{-1}(\mathbb{1}).$$

Now since A is a subHopf algebra, we have

$$\begin{aligned} (Z \otimes id)\Delta(T(x)) &\subseteq A[[x]] \\ \Rightarrow (id - L)^{-1}(\mathbb{1}) &\subseteq A[[x]]. \end{aligned}$$

Now the third step is to pull out easy coefficient and compare. From $T(x) = xB_+(P(T(x)))$ we have the recursive expression

$$t_1 = \bullet, t_{n+1} = \sum_{k=1}^n \sum_{\alpha_1 + \dots + \alpha_k = n} p_k B_+(t_{\alpha_1} \dots t_{\alpha_k}).$$

Write $(id - L)^{-1}(\mathbb{1}) = \sum_{i=0}^{\infty} b_i x^i$. By induction

$$b_0 = 1$$

$$\begin{aligned} b_{n+1} &= \sum_{k=1}^n \sum_{\alpha_1 + \dots + \alpha_k = n} (k+1)p_{k+1} B_+(t_{\alpha_1} \dots t_{\alpha_k}) \\ &\quad + \sum_{k=1}^n \sum_{\alpha_1 + \dots + \alpha_k = n} k p_k B_+(b_{\alpha_1} t_{\alpha_2} \dots t_{\alpha_k}). \end{aligned}$$

Now compare coefficients. Consider $B_+(f_n), B_+(b_n)$, i.e. trees where the root has only one child and degree $n+1$ in f_{n+1} and b_{n+1} . Coefficient in f_{n+1} is $p_1 B_+(f_n)$, and the coefficient in b_{n+1} is $2p_2 B_+(f_n) + p_1 B_+(b_n)$, but by assumption the f_n make a subHopf algebra, and the b_n are in it, so $b_{n+1} = \lambda_{n+1} f_{n+1}, \lambda_{n+1} \in \mathbb{Q}$. So

$$\lambda_1 = p_1, \lambda_{n+1} = \left(\frac{2p_2}{p_1} + \lambda_n \right). \quad (2)$$

Consider $B_y(\bullet^n)$ in f_{n+1} and b_{n+1} . In f_{n+1} we get p_n , and in b_{n+1} we get $(n+1)p_{n+1} + np_n p_1$, so

$$\lambda_{n+1} p_n = (n+1)p_{n+1} + np_n p_1, \quad \forall n \geq 1,$$

which together with (2) gives

$$(n+1)p_{n+1} + (p_1 - 2\frac{p_2}{p_1})np_n = p_1 p_n.$$

If we rewrite this at level of series we get

$$P'(h) + (p_1 - 2\frac{p_2}{p_1})hP'(h) = p_1 P(h).$$

Now let $\alpha = p_1, \beta = 2\frac{p_2}{p_1} - 1$ to get the result.

△

References

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